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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
10/056,609	01/24/2002	Ronald L. Pettyjohn	FORE-609	9506
7590	03/28/2006		EXAMINER	
Ansel M. Schwartz Attorney at Law Suite 304 201 N. Craig Street Pittsburgh, PA 15222			WALSH, JOHN B	
			ART UNIT	PAPER NUMBER
			2151	
			DATE MAILED: 03/28/2006	

Please find below and/or attached an Office communication concerning this application or proceeding.

Office Action Summary	Application No.	Applicant(s)	
	10/056,609	PETTYJOHN ET AL.	
	Examiner John B. Walsh	Art Unit 2151	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

- 1) Responsive to communication(s) filed on 03 January 2006.
- 2a) This action is **FINAL**. 2b) This action is non-final.
- 3) Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

- 4) Claim(s) 1-20 is/are pending in the application.
 - 4a) Of the above claim(s) _____ is/are withdrawn from consideration.
- 5) Claim(s) _____ is/are allowed.
- 6) Claim(s) 1-20 is/are rejected.
- 7) Claim(s) _____ is/are objected to.
- 8) Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

- 9) The specification is objected to by the Examiner.
- 10) The drawing(s) filed on _____ is/are: a) accepted or b) objected to by the Examiner.

Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).

Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

- 12) Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
 - a) All b) Some * c) None of:
 1. Certified copies of the priority documents have been received.
 2. Certified copies of the priority documents have been received in Application No. _____.
 3. Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

* See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

1) <input type="checkbox"/> Notice of References Cited (PTO-892)	4) <input type="checkbox"/> Interview Summary (PTO-413)
2) <input type="checkbox"/> Notice of Draftsperson's Patent Drawing Review (PTO-948)	Paper No(s)/Mail Date. _____
3) <input type="checkbox"/> Information Disclosure Statement(s) (PTO-1449 or PTO/SB/08) Paper No(s)/Mail Date _____	5) <input type="checkbox"/> Notice of Informal Patent Application (PTO-152)
	6) <input type="checkbox"/> Other: _____

DETAILED ACTION

Claim Rejections - 35 USC § 102

1. The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless –

(b) the invention was patented or described in a printed publication in this or a foreign country or in public use or on sale in this country, more than one year prior to the date of application for patent in the United States.
(e) the invention was described in (1) an application for patent, published under section 122(b), by another filed in the United States before the invention by the applicant for patent or (2) a patent granted on an application for patent by another filed in the United States before the invention by the applicant for patent, except that an international application filed under the treaty defined in section 351(a) shall have the effects for purposes of this subsection of an application filed in the United States only if the international application designated the United States and was published under Article 21(2) of such treaty in the English language.

2. Claims 1-5 and 7 are rejected under 35 U.S.C. 102(b) as being anticipated by U.S. Patent No. 5,940,397 to Gritton.

As concerns claim 1, a method of buffering at least one data unit received at a node on a communications network, the at least one data unit being associated with one or more channels in the network, the network node including at least one input port communicably coupleable to at least one output port, comprising the steps of: providing a memory (46) at each output port of the network node, the memory comprising at least one linear time-indexed array having a plurality of locations for buffering (column 7, lines 30-40) the at least one data unit; in the event the at least one data unit is associated with a single channel in the network, storing the at least one data unit in a respective location of the time-indexed array (column 7, lines 30-50); and in the event the at least one data unit is associated with a plurality of channels in the network, dividing the memory (42, column 11, line 1) into a plurality of linear time-indexed arrays, each time-indexed array corresponding to a respective channel, and storing the at least one data unit in a respective location of the corresponding time-indexed array (regardless of the

number of channels, such that if there is one channel or a plurality of channels it will perform the functions, column 7, lines 30-50).

As concerns claim 2, the method of claim 1 wherein the network node has predetermined total bandwidth (inherent for any particular device to have a physical limit for bandwidth, column 3, lines 19-22), and the providing step includes providing a memory within the node, the memory having a size sufficient to support the total bandwidth of the node.

As concerns claim 3, the method of claim 1 wherein the dividing step includes dividing the memory into a plurality of arrays (arrays in memory, column 7, lines 23-24), each array corresponding to a respective channel (corresponds to data received over channels), the respective channels conforming to predetermined bandwidth requirements (inherent for any particular device to have a physical limit for bandwidth).

As concerns claim 4, the method of claim 3 wherein the dividing step includes dividing the memory into a plurality of arrays (arrays in memory, column 7, lines 23-24), each array having a size proportional to a fractional amount of a predetermined total bandwidth of the node (column 7, lines 38-40, inherent that the amount of data stored is proportional to the bandwidth).

As concerns claim 5, the method of claim 1 wherein the at least one data unit has an associated timestamp value and the first storing step includes storing the at least one data unit in a respective location of the time-based array based on the associated timestamp value (column 7, line 54, column 12, lines 34-36).

As concerns claim 7, the method of claim 1 wherein the at least one data unit has an associated timestamp value, and the second storing step includes storing the at least one data unit in a respective location of the corresponding array based on the associated timestamp value (column 7, line 54, column 12, lines 34-36).

3. Claims 1-20 are rejected under 35 U.S.C. 102(e) as being anticipated by U.S. Patent No. 6,081,507 to Chao et al.

As concerns claim 1, a method of buffering at least one data unit received at a node on a communications network, the at least one data unit being associated with one or more channels in the network, the network node including at least one input port communicably coupleable to at least one output port, comprising the steps of: providing a memory (1600) at each output port of the network node, the memory comprising at least one linear time-indexed array having a plurality of locations for buffering (column 18, lines 15-16, column 16, line 16) the at least one data unit; in the event the at least one data unit is associated with a single channel in the network, storing the at least one data unit in a respective location of the time-indexed array (column 15, lines 24-25); and in the event the at least one data unit is associated with a plurality of channels in the network, dividing the memory (column 13, lines 18-19) into a plurality of linear time-indexed arrays, each time-indexed array corresponding to a respective channel, and storing the at least one data unit in a respective location of the corresponding time-indexed array (regardless of the number of channels, such that if there is one channel or a plurality of channels it will perform the functions).

As concerns claim 2, the method of claim 1 wherein the network node has predetermined total bandwidth (inherent for any particular device to have a physical limit for bandwidth, column 13, line 15), and the providing step includes providing a memory within the node, the memory having a size sufficient to support the total bandwidth of the node.

As concerns claim 3, the method of claim 1 wherein the dividing step includes dividing the memory into a plurality of arrays (arrays in memory,), each array corresponding to a respective channel (corresponds to data received over channels, column 13, line 15), the respective channels conforming to predetermined bandwidth requirements (inherent for any particular device to have a physical limit for bandwidth).

As concerns claim 4, the method of claim 3 wherein the dividing step includes dividing the memory into a plurality of arrays (arrays in memory, column 16, line 16), each array having a size proportional to a fractional amount of a predetermined total bandwidth of the node (inherent that the amount of data stored is proportional to the bandwidth).

As concerns claim 5, the method of claim 1 wherein the at least one data unit has an associated timestamp value and the first storing step includes storing the at least one data unit in a respective location of the time-based array based on the associated timestamp value (column 16, line 16).

As concerns claims 6 and 8, weighted-fair queuing algorithm (column 14, line 58).

As concerns claim 7, the method of claim 1 wherein the at least one data unit has an associated timestamp value, and the second storing step includes storing the at least one data unit in a respective location of the corresponding array based on the associated timestamp value (column 16, line 16).

As concerns claim 9, a method of scheduling the transmission of at least one data unit from a node on a communications network, the network node including at least one input port communicably coupleable to at least one output port, comprising the steps of: providing a first memory (1600) at each output port (column 8, line 21) of the network node, the first memory comprising at least one linear time-indexed array having a plurality of locations for buffering

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the at least one data unit (column 18, lines 15-16; column 16, line 16); receiving at least one first data unit at the network node, the at least one first data unit having an associated timestamp value (column 16, line 16); inserting the first data unit into a respective location of the time-based array based on the associated timestamp value (column 18, line 16); partitioning a binary value of the timestamp associated with the first data unit into a plurality of sub-fields, each sub-field comprising one or more bits, and using the plurality of sub-fields to generate a corresponding plurality of acceleration bit-strings for use in identifying the first data unit in the time-based array having a lowest associated timestamp value (figures 19A, 19B, 20A, 20B); extracting the identified first data unit having the lowest associated timestamp value from the time-based array (figures 19A, 19B, 20A, 20B); and designating the extracted first data unit as a next data unit to be transmitted over the network (figure 28, 2840).

As concerns claim 10, the method of claim 9 wherein the network node includes at least one second memory (1600, a second memory location or address) and the partitioning step includes using the sub-fields of bits to index respective locations in the at least one second memory and asserting values at the respective memory locations to generate the plurality of acceleration bit-strings.

As concerns claim 11, the method of claim 10 further including the step of priority encoding each acceleration bit-string to obtain a corresponding priority-encoded acceleration bit-string (figures 19A, 19B, 20A, 20B; column 17, lines 15-35).

As concerns claim 12, the method of claim 11 wherein the priority encoding step employs "low-wins" priority encoding (column 15, line 25).

As concerns claim 13, the method of claim 11 further including the step of employing one or more of the priority-encoded acceleration bit-strings to index the time-based array to

identify the first data unit in the array having the lowest associated timestamp value (figures 19A, 19B, 20A, 20B, column 17, lines 15-35, column 18, lines 15-16).

As concerns claim 14, the method of claim 9 wherein each location of the time-based array corresponds to a respective timestamp value within a first time window ranging from $t=0$ to $t=Tw$, and further including the step of in the event the timestamp value associated with the next data unit to be transmitted over the network is greater than or equal to $Tw/2$, shifting the first time window forward in time by $Tw/2$ to obtain a next time window ranging from $t=Tw/2$ to $t=3Tw/2$ (column 15, lines 24-25, figures 34a, 34b).

As concerns claim 15, the method of claim 14 wherein the receiving step includes receiving at least one first data unit at the network node, the at least one first data unit having an associated timestamp value within a range limited to $Tw/2$ (column 15, lines 24-25, figures 34a, 34b).

As concerns claim 16, a system for scheduling the transmission of at least one data unit from a node on a communications network, the node including at least one input port and at least one output port, the input port being communicably coupleable to the output port, comprising: a first memory (1600) disposed at each output port of the network node, the first memory comprising a linear time-indexed array having a plurality of locations configured to buffer at least one first data unit, each first data unit having an associated timestamp value (column 18, line 16); and a controller configured to insert the at least one first data unit into a respective location of the time-based array based on the associated timestamp value, partition a binary value of the timestamp associated with the first data unit into a plurality of sub-fields, each sub-field comprising one or more bits, use the plurality of sub-fields to generate a

corresponding plurality of acceleration bit-strings for use in identifying the first data unit in the time-based array having a lowest associated timestamp value, extract the identified first data unit having the lowest associated timestamp value from the time-based array, and designate the extracted first data unit as a next data unit to be transmitted over the network (figures 19A, 19B, 20A, 20B, figure 28, 2840).

As concerns claim 17, the system of claim 16 wherein each location of the time-based array corresponds to a respective timestamp value within a first time window ranging from $t=0$ to $t=Tw$, and the controller is further configured to, in the event the timestamp value associated with the next data unit to be transmitted over the network is greater than or equal to $Tw/2$, shift the first time window forward in time by $Tw/2$ to obtain a next time window ranging from $t=Tw/2$ to $t=3Tw/2$ (column 15, lines 24-25, figures 34a, 34b).

As concerns claim 18, the system of claim 16 wherein the network node has a predetermined total bandwidth and the time-based array has a size sufficient to support the total bandwidth of the node (over time it can support the total bandwidth).

As concerns claim 19, the system of claim 16 wherein the first memory comprises a plurality of linear time-indexed arrays, each array corresponding to a respective channel in the network (column 13, lines 18-19).

As concerns claim 20, the system of claim 19 wherein each array has a size proportional to a fractional amount of a predetermined total bandwidth of the network node (inherent that the amount of data stored is proportional to the bandwidth).

Claim Rejections - 35 USC § 103

4. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

5. Claims 6 and 8 are rejected under 35 U.S.C. 103(a) as being unpatentable over U.S.

Patent No. 5,940,397 to Gritton as applied to claims 5 and 7 above in view of U.S. Patent No. 6,081,507 to Chao et al.

Gritton '397 does not explicitly disclose a weighted-fair queuing algorithm.

Chao et al. '507 teach a weighted-fair queuing algorithm (column 14, line 58).

It would have been obvious to one having ordinary skill in the art at the time the invention was made to provide Gritton '397 with a weighted-fair queuing algorithm, as taught by Chao et al. '507, in order to provide improved control and traffic management.

Response to Arguments

6. Applicant's arguments filed January 3, 2006 have been fully considered but they are not persuasive.

The applicant argues Gritton does not disclose a linear time-indexed array. The examiner disagrees since Gritton discloses a linear time-indexed array (figure 6-linear array; column 6, lines 34 and 39-42 – sequentially indexing and the list is dependent upon a time value).

The applicant argues Chao does not disclose “in the event the at least one data unit is associated with a plurality of channels in the network, dividing the memory into a plurality of liner time-indexed arrays, each time-indexed array corresponding to a respective channel”. The

examiner disagrees since Chao discloses having data units associated with a plurality of channels and dividing the memory into time-indexed arrays corresponding to respective channels (column 13, lines 18-19- a separate queue may be provided for each flow). In response to applicant's argument that the references fail to show certain features of applicant's invention, it is noted that the features upon which applicant relies (i.e., a data unit associated with a plurality of channels; the claims recite at least one data unit, implying a plurality of units not just a singular data unit) are not recited in the rejected claim(s). Although the claims are interpreted in light of the specification, limitations from the specification are not read into the claims. See *In re Van Geuns*, 988 F.2d 1181, 26 USPQ2d 1057 (Fed. Cir. 1993).

Conclusion

7. **THIS ACTION IS MADE FINAL.** Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

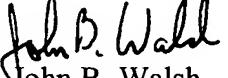
A shortened statutory period for reply to this final action is set to expire THREE MONTHS from the mailing date of this action. In the event a first reply is filed within TWO MONTHS of the mailing date of this final action and the advisory action is not mailed until after the end of the THREE-MONTH shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than SIX MONTHS from the mailing date of this final action.

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8. Any inquiry concerning this communication or earlier communications from the examiner should be directed to John B. Walsh whose telephone number is 571-272-7063. The examiner can normally be reached on Monday-Wednesday from 5:30-4:00.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Zarni Maung can be reached on 571-272-3939. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free).


John B. Walsh
Primary Examiner
Art Unit 2151